

Rate-Distortion Analysis for Sampled Correlated Cyclostationary Gaussian Processes

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ECE Departmental Seminar

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 - Cyclostationary Processes
 - Statistics of Sampled CT WSCS Processes
 - Rate-Distortion Theory
 - Information-Spectrum Framework
- 2 Motivating Example and Related Works
 - Compress-and-Forward Relay Networks
 - Related Works
- 3 Problem Formulation
 - Definitions of Related Information-Spectrum Quantities
 - Source Sequence Generation Model
 - The RDF for Compressing an Arbitrary DT Process
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 - Proof Sketch
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 - Auxiliary Theorem: The RDF for DT WSCS Processes
 - The RDF with a Finite and Bounded Delay between Sequences
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 - Numerical Evaluations and Implications
- 6 Summary and Future Works

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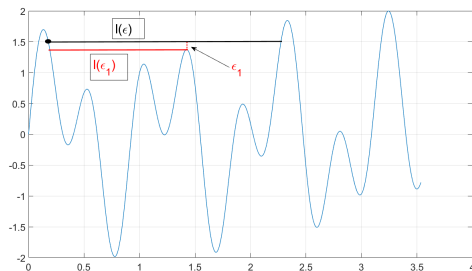
- Man-made signals are typically generated via **repetitive** operations, resulting in **cyclostationary** statistics.
 - A **wide-sense cyclostationary (WSCS)** process has a *periodic mean* and a *periodic autocorrelation function (AF)* with the *same* period:
 - A real **continuous-time (CT)** random process $X(t)$, $t \in \mathcal{R}$, is called **WSCS** if $\exists \mathbf{T}_c \in \mathcal{R}^{++}$, s.t. $\forall \lambda \in \mathcal{R}$,

$$m_X(t) \triangleq \mathbb{E}\{X(t)\} = m_X(t + \mathbf{T}_c),$$
$$c_X(t, \lambda) \triangleq \mathbb{E}\{X(t) \cdot X(t + \lambda)\} = c_X(t + \mathbf{T}_c, \lambda).$$

- Similarly, we define a **discrete-time (DT) WSCS** process.
- For example, *baseband OFDM signals* are **CT WSCS**.

- A real DT deterministic function $f[i]$, $i \in \mathcal{Z}$, is called **almost periodic**, if
 - $\forall \epsilon \in \mathcal{R}^{++}$, $\exists l_\epsilon \in \mathcal{N}^+$, which satisfies:
 - $\forall \alpha \in \mathcal{Z}$, $\exists \Delta \in [\alpha, \alpha + l_\epsilon)$, s.t.

$$\sup_{i \in \mathcal{Z}} |f[i + \Delta] - f[i]| < \epsilon.$$



- A real **zero-mean** DT random process $X[i]$, $i \in \mathcal{Z}$, is called **wide-sense almost cyclostationary (WSACS)** if $c_X[i, \Delta]$ is **almost periodic** in $i \in \mathcal{Z}$.

Sampling a CT WSCS process

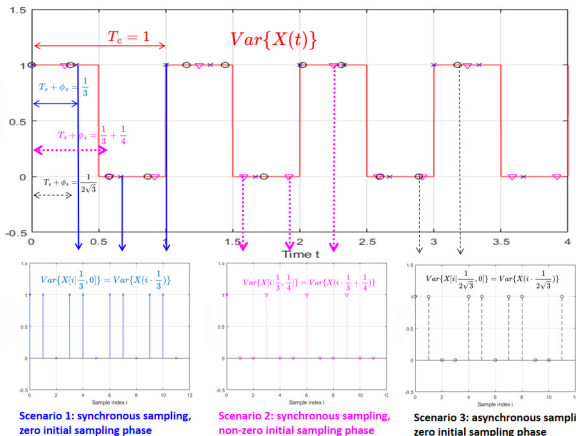
- Given a **CT WSCS source process** $X_c(t)$ with the **period of statistics** of T_c .
- Sample $X_c(t)$ using the **interval** $T_s(\epsilon) \triangleq \frac{T_c}{p+\epsilon}$, where $p \in \mathcal{N}^+$ and $\epsilon \in (0, 1)$, and the **initial sampling phase** $\phi_s \in [0, T_c)$

→ DT sampled process:

$$X_\epsilon^{\phi_s}[j] \triangleq X_c(T_s(\epsilon) \cdot j + \phi_s).$$

- Synchronous sampling:** $\epsilon \in \mathcal{Q}$ → $X_\epsilon^{\phi_s}[j]$ is a **DT WSCS process**;
- Asynchronous sampling:** $\epsilon \notin \mathcal{Q}$ → $X_\epsilon^{\phi_s}[j]$ is a **DT WSACS process**.

Sampling a CT WSCS process



Sampling a **CT WSCS** process vs. Sampling a **CT stationary** process

NOT necessarily a **DT WSCS** process. **Necessarily** a **DT stationary** process.

- Denote with \mathcal{X} the **source alphabet** and with $\hat{\mathcal{X}}$ the **reconstruction alphabet**.

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- A **lossy source code** (M, l) consists of
 - a **message set** of **size** M ;
 - an **encoder**: $f_l(\cdot) : \mathcal{X}^l \mapsto \{0, 1, 2, \dots, M - 1\}$;
 - a **decoder**: $g_l(\cdot) : \{0, 1, 2, \dots, M - 1\} \mapsto \hat{\mathcal{X}}^l$.

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- The **squared-error distortion function** between a **source sequence** $\{x_i\}_{i=0}^{l-1} \equiv x^l$ and a **reconstruction sequence** $\{\hat{x}_i\}_{i=0}^{l-1} \equiv \hat{x}^l$ is given as

$$d_{se}(x^l, \hat{x}^l) \triangleq \frac{1}{l} \sum_{i=0}^{l-1} (x_i - \hat{x}_i)^2 \equiv \bar{d}(x^l, \hat{x}^l).$$

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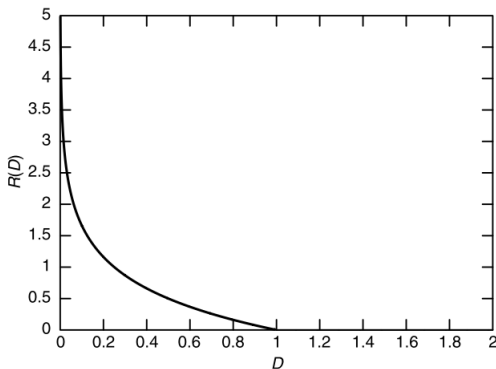
- A **rate-distortion pair** (R, D) is **achievable**, if there exists a **lossy source code** $(2^{lR}, l)$ satisfying

$$\limsup_{l \rightarrow \infty} \mathbb{E} \left\{ \bar{d}(\mathbf{X}^l, g_l(f_l(\mathbf{X}^l))) \right\} \leq D.$$

- The **rate-distortion function (RDF)** is the **infimum** of all **code rates** R , for which the **rate-distortion pair** (R, D) is **achievable**.

- **RDF** for an **i.i.d. Gaussian source** $\mathcal{N}(0, \sigma^2)$:

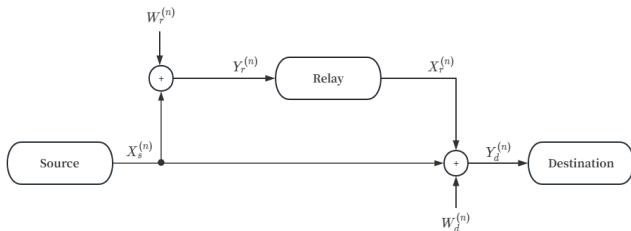
$$R(D) = \begin{cases} \frac{1}{2} \log \frac{\sigma^2}{D}, & 0 \leq D \leq \sigma^2, \\ 0, & D > \sigma^2. \end{cases}$$



- **Standard rate-distortion analysis** relies on the **stationarity** and the **ergodicity** (or the **information-stability**) of the source.
- However, DT WSACS processes are **nonstationary** and **nonergodic**.
⇒ **Conventional** information-theoretic arguments are **inapplicable**.
- Suitable frameworks for analyzing **information-unstable processes**:
 - **Asymptotically mean-stationary (AMS) processes** [Gray: 2011];
 - **Information-spectrum framework** [Han: 2010].
- We use the **information-spectrum framework** in our study.

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Compress-and-forward relay networks

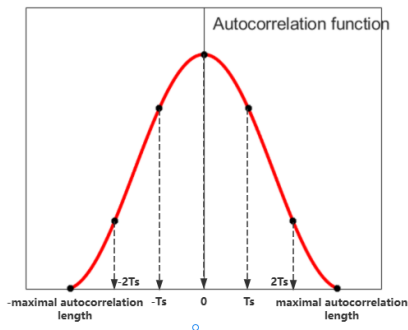


- The (*Gaussian*) signal received at the **relay** is **sampled**,
 - before being **compressed** and **forwarded** to the destination.
- Practically, the **ratio** between the **actual sampling interval** at the relay and the **actual period of statistics** of the received signal is **irrational**, due to
 - **physical separation** of the **relay oscillator** and the **source oscillator**;
 - **inherent variability** of oscillators' frequencies (i.e., **jitter**).

Motivating Example and Related Works

Compress-and-Forward Relay Networks

- To reduce the **loss of information**, the **sampling interval** at the relay is typically *smaller* than the **maximal autocorrelation length** of the received signal.



Therefore, the sampled received signal at the relay is a **DT WSACS Gaussian process with memory**.

- [Kipnis et al.: *IEEE TIT* 2018] characterized the **distortion-rate function (DRF)** of **DT WSCS Gaussian processes**,
 - by transforming the **DT scalar WSCS process** into a **DT vector stationary process**.

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 - using the **information-spectrum framework**.
- The **capacity** of **additive DT WSACS Gaussian channels** was characterized using the **information-spectrum framework**,
 - in [Shlezinger et al.: *IEEE TCOM* 2020] for memoryless noise;
 - in [Dabora and Abakasanga: *IEEE TIT* 2023] for noise with memory.

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Our task: RDF characterization for compressing DT WSACS Gaussian processes with memory.

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- Consider a sequence of real random variables (RVs) $\{X_i\}_{i=0}^{\infty}$:
 - The **limit superior in probability** is given as

$$p\text{-}\limsup_{i \rightarrow \infty} X_i \triangleq \inf \{ \alpha \in \mathcal{R} \mid \lim_{i \rightarrow \infty} \Pr\{X_i > \alpha\} = 0 \};$$

- Given a **common probability measure** P , if

$$\lim_{u \rightarrow \infty} \sup_{i \geq 0} \int_{|X_i| \geq u} |X_i| dP = 0,$$

then the sequence is said to be **uniformly integrable**.

Problem Formulation

Definitions of Related Information-Spectrum Quantities

Consider a real DT **source process** $X[i]$, $i \in \mathcal{N}$, and a **reconstruction process** $\hat{X}[i]$:

- Denote a **block of l symbols** $\{X[i]\}_{i=0}^{l-1} \equiv \mathbf{X}^l$;
- Denote a **block of l reconstruction symbols** $\{\hat{X}[i]\}_{i=0}^{l-1} \equiv \hat{\mathbf{X}}^l$.
- Let $f_{\hat{\mathbf{X}}^l | \mathbf{X}^l}(\hat{\mathbf{x}}^l | \mathbf{x}^l)$ denote the **conditional probability density function (PDF)** of $\hat{\mathbf{X}}^l$ given \mathbf{X}^l .
- Let $f_{\hat{\mathbf{X}}^l}(\hat{\mathbf{x}}^l)$ denote the **PDF** of $\hat{\mathbf{X}}^l$.

Problem Formulation

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- Let $f_{\hat{\mathbf{X}}^l|\mathbf{X}^l}(\hat{\mathbf{x}}^l|\mathbf{x}^l)$ denote the **conditional PDF** of $\hat{\mathbf{X}}^l$ given \mathbf{X}^l .
- Let $f_{\hat{\mathbf{X}}^l}(\hat{\mathbf{x}}^l)$ denote the **PDF** of $\hat{\mathbf{X}}^l$.
- The **mutual information density rate** between \mathbf{X}^l and $\hat{\mathbf{X}}^l$ is defined as

$$i(\mathbf{X}^l; \hat{\mathbf{X}}^l) \triangleq \frac{1}{l} \log \frac{f_{\hat{\mathbf{X}}^l|\mathbf{X}^l}(\hat{\mathbf{x}}^l|\mathbf{x}^l)}{f_{\hat{\mathbf{X}}^l}(\hat{\mathbf{x}}^l)}.$$

- The **spectral sup-mutual information rate** between \mathbf{X}^∞ and $\hat{\mathbf{X}}^\infty$, $\bar{I}(\mathbf{X}^\infty, \hat{\mathbf{X}}^\infty)$, is defined as

$$\bar{I}(\mathbf{X}^\infty, \hat{\mathbf{X}}^\infty) \triangleq \text{p-} \limsup_{l \rightarrow \infty} i(\mathbf{X}^l; \hat{\mathbf{X}}^l).$$

- Consider a real **CT WSCS Gaussian** source process $X_C(t)$, s.t.
 - $\mathbb{E}\{X_C(t)\} = 0$.
 - AF $c_{X_C}(t, \lambda)$ is **periodic** in t with **period** T_C , **bounded** and **uniformly continuous** in both t and λ .
 - AF has **finite-memory** with maximal autocorrelation length λ_C ,
 - ▶ i.e., $c_{X_C}(t, \lambda) \triangleq \mathbb{E}\{X_C(t) \cdot X_C(t + \lambda)\} = 0, \forall |\lambda| > \lambda_C$.

Problem Formulation

Source Sequence Generation Model

- Consider a real **CT WSCS Gaussian** source process $X_c(t)$, s.t.
 - $\mathbb{E}\{X_c(t)\} = 0$.
 - AF $c_{X_c}(t, \lambda)$ is **periodic** in t with **period** T_c , **bounded** and **uniformly continuous** in both t and λ .
 - AF has **finite-memory** with maximal autocorrelation length λ_c ,
 - ▶ i.e., $c_{X_c}(t, \lambda) \triangleq \mathbb{E}\{X_c(t) \cdot X_c(t + \lambda)\} = 0, \forall |\lambda| > \lambda_c$.
- $X_c(t)$ is **sampled** with **sampling interval** $T_s(\epsilon) \leq \lambda_c$ and **initial sampling phase** $\phi_s \in [0, T_c) \rightarrow$ DT sampled process:

$$X_\epsilon^{\phi_s}[i] \triangleq X_c(T_s(\epsilon) \cdot i + \phi_s),$$

$T_s(\epsilon) \triangleq \frac{T_c}{p+\epsilon}$, where $p \in \mathcal{N}^+$ and $\epsilon \notin \mathcal{Q}$, $\epsilon \in (0, 1)$ is the **sampling mismatch**.

- Consider a real **CT WSCS Gaussian** source process $X_c(t)$, s.t.
 - $\mathbb{E}\{X_c(t)\} = 0$.
 - AF $c_{X_c}(t, \lambda)$ is **periodic** in t with **period** T_c , **bounded** and **uniformly continuous** in both t and λ .
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- $X_c(t)$ is **sampled** with **sampling interval** $T_s(\epsilon) \leq \lambda_c$ and **initial sampling phase** $\phi_s \in [0, T_c) \rightarrow$ DT sampled process:

$$X_\epsilon^{\phi_s}[i] \triangleq X_c(T_s(\epsilon) \cdot i + \phi_s),$$

$T_s(\epsilon) \triangleq \frac{T_c}{p+\epsilon}$, where $p \in \mathcal{N}^+$ and $\epsilon \notin \mathcal{Q}$, $\epsilon \in (0, 1)$ is the **sampling mismatch**.

- Source sequences can be generated
 - ① **continuously**, i.e., delay between consecutive sequences is **NOT** allowed;
 - ② with **finite and bounded delay** between consecutive sequences.
 - \Rightarrow all sequences are **statistically independent**.

Problem Formulation

The RDF for Compressing an Arbitrary DT Process

- For an **arbitrary** DT process $X[i]$, $i \in \mathcal{N}$, if for a **block of l symbols** $\{X[i]\}_{i=0}^{l-1} \equiv \mathbf{X}^l$, there exists a **deterministic reference word** $\{r_i\}_{i=0}^{l-1} \equiv \mathbf{r}^l$ which makes the **sequence of RVs** $\{\bar{d}(\mathbf{X}^l, \mathbf{r}^l)\}_{l=0}^{\infty}$ **uniformly integrable**.
 - Denote the **block of l reconstruction symbols** $\{\hat{X}_i\}_{i=0}^{l-1} \equiv \hat{\mathbf{X}}^l$ and the **joint cumulative distribution function (CDF)** of \mathbf{X}^l and $\hat{\mathbf{X}}^l$ as $F_{\mathbf{X}^l, \hat{\mathbf{X}}^l}$.
 - Then, the **RDF** for compressing $X[i]$ using the **fixed-length coding** and the **average distortion criterion** is [Han: 2010]

$$R(D) = \limsup_{l \rightarrow \infty} \inf_{F_{\mathbf{X}^l, \hat{\mathbf{X}}^l}: \mathbb{E}\{\bar{d}(\mathbf{X}^l, \hat{\mathbf{X}}^l)\} \leq D} \bar{I}(\mathbf{X}^l, \hat{\mathbf{X}}^l).$$

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Main Theorem 1

The RDF for Continuous Source Sequence Generation

- The **RDF for compressing DT WSACS Gaussian processes with memory with continuous source sequence generation**:

- Denote a **block of l symbols** $\{X_{\epsilon}^{\phi_s}[i]\}_{i=0}^{l-1} \equiv \mathbf{X}_{\epsilon, \phi_s}^l$, its **block of l reconstruction symbols** $\{\hat{X}_{\epsilon}^{\phi_s}[i]\}_{i=0}^{l-1} \equiv \hat{\mathbf{X}}_{\epsilon, \phi_s}^l$, and

$$\mathbf{S}_{\epsilon, \phi_s}^l \triangleq \mathbf{X}_{\epsilon, \phi_s}^l - \hat{\mathbf{X}}_{\epsilon, \phi_s}^l.$$

- Define the set

$$\mathcal{C}_{\mathbf{S}_{\epsilon, \phi_s}^l} \triangleq \left\{ \mathbf{C}_{\mathbf{S}_{\epsilon, \phi_s}^l} \in \mathcal{R}^{l \times l} \left| \frac{1}{l} \text{tr}\{\mathbf{C}_{\mathbf{S}_{\epsilon, \phi_s}^l}\} \leq D, \mathbf{C}_{\mathbf{S}_{\epsilon, \phi_s}^l} \succcurlyeq 0, \right. \right. \\ \left. \left. \mathbf{C}_{\mathbf{S}_{\epsilon, \phi_s}^l} = (\mathbf{C}_{\mathbf{S}_{\epsilon, \phi_s}^l})^T, \mathbf{C}_{\mathbf{X}_{\epsilon, \phi_s}^l} \succeq \mathbf{C}_{\mathbf{S}_{\epsilon, \phi_s}^l} \right\}.$$

Main Theorem 1

The RDF for Continuous Source Sequence Generation

- Then the **RDF** $R_\epsilon(D)$ for compressing $X_\epsilon^{\phi_s}[i]$ is

$$R_\epsilon(D) = \frac{1}{T_c} \int_{\phi_s=0}^{T_c} R_\epsilon^{\phi_s}(D) d\phi_s,$$

where

$$R_\epsilon^{\phi_s}(D) \triangleq \limsup_{l \rightarrow \infty} \min_{\mathbf{C}_{S^l}^{\epsilon, \phi_s} \in \mathcal{C}_{S^l}^{\epsilon, \phi_s}} \frac{1}{2l} \log \left(\frac{\det(\mathbf{C}_{\mathbf{X}^l}^{\epsilon, \phi_s})}{\det(\mathbf{C}_{S^l}^{\epsilon, \phi_s})} \right),$$

in which $\mathbf{S}_{\epsilon, \phi_s}^l$ is a **Gaussian vector**.

- **Corollary (RDF with the fixed code rate)**

- If the **code rate** has to be **fixed**, the **RDF** is

$$R_\epsilon(D) = \min_{\phi_s \in [0, T_c)} R_\epsilon^{\phi_s}(D).$$

Converse Part:

- 1 Consider a *specific sequence of lossy source codes* $(\tilde{M}(l, D), l)$ defined by the **encoder** $\tilde{f}_l(\cdot)$ and the **decoder** $\tilde{g}_l(\cdot)$, such that

$$\limsup_{l \rightarrow \infty} \mathbb{E}\{\bar{d}_{se}(\mathbf{X}_{\epsilon, \phi_s}^l, \hat{\mathbf{X}}_{\epsilon, \phi_s}^l)\} \leq D,$$

where $\hat{\mathbf{X}}_{\epsilon, \phi_s}^l = \tilde{g}_l(\tilde{f}_l(\mathbf{X}_{\epsilon, \phi_s}^l))$.

- 2 By [Kostina2013], we obtain

$$\log \tilde{M}(l, D) \geq \inf_{f(\hat{\mathbf{X}}_{\epsilon, \phi_s}^l | \mathbf{X}_{\epsilon, \phi_s}^l): \mathbb{E}\{\bar{d}_{se}(\mathbf{X}_{\epsilon, \phi_s}^l, \hat{\mathbf{X}}_{\epsilon, \phi_s}^l)\} \leq \mathbb{E}\{\bar{d}_{se}(\mathbf{X}_{\epsilon, \phi_s}^l, \hat{\mathbf{X}}_{\epsilon, \phi_s}^l)\}} I(\mathbf{X}_{\epsilon, \phi_s}^l; \hat{\mathbf{X}}_{\epsilon, \phi_s}^l).$$

Main Theorem 1

Proof Sketch

- By **asymptotic properties** of the **limit superior** and the condition $\limsup_{l \rightarrow \infty} \mathbb{E}\{\bar{d}_{se}(\mathbf{X}'_{\epsilon, \phi_s}, \hat{\mathbf{X}}'_{\epsilon, \phi_s})\} \leq D$, we obtain

$$R_{\epsilon}^{\phi_s}(D) \geq \limsup_{l \rightarrow \infty} \inf_{\substack{f(\hat{\mathbf{X}}'_{\epsilon, \phi_s} | \mathbf{X}'_{\epsilon, \phi_s}): \\ \mathbb{E}\{\bar{d}_{se}(\mathbf{X}'_{\epsilon, \phi_s}, \hat{\mathbf{X}}'_{\epsilon, \phi_s})\} \leq D}} \frac{1}{l} I(\mathbf{X}'_{\epsilon, \phi_s}; \hat{\mathbf{X}}'_{\epsilon, \phi_s}).$$

- Minimizing $\frac{1}{l} I(\mathbf{X}'_{\epsilon, \phi_s}; \hat{\mathbf{X}}'_{\epsilon, \phi_s})$, the **optimal** $\mathbf{S}'_{\epsilon, \phi_s}$ is found to be **Gaussian**, therefore

$$R_{\epsilon}^{\phi_s}(D) \geq \limsup_{l \rightarrow \infty} \min_{\mathbf{C}_{\mathbf{S}'_{\epsilon, \phi_s}} \in \mathcal{C}_{\mathbf{S}'_{\epsilon, \phi_s}}} \frac{1}{2l} \log \left(\frac{\det(\mathbf{C}_{\mathbf{X}'_{\epsilon, \phi_s}})}{\det(\mathbf{C}_{\mathbf{S}'_{\epsilon, \phi_s}})} \right).$$

Achievability Part:

- 1 Denote the **optimal block of l reconstruction symbols** as $\{\hat{\mathbf{X}}_{\epsilon, \phi_s}^{l,*}\}_{i=0}^{l-1} \equiv \hat{\mathbf{X}}_{\epsilon, \phi_s}^{l,*}$ and the **mutual information density rate** between $\mathbf{X}_{\epsilon, \phi_s}^l$ and $\hat{\mathbf{X}}_{\epsilon, \phi_s}^{l,*}$ as $Z_l(F_{X_{\epsilon, \hat{X}_{\epsilon}^*}}^{\phi_s})$.

The **mutual information density** $V_{\epsilon, l}^{\phi_s}$ between $\mathbf{X}_{\epsilon, \phi_s}^l$ and $\hat{\mathbf{X}}_{\epsilon, \phi_s}^{l,*}$ is

$$V_{\epsilon, l}^{\phi_s} = l \cdot \underbrace{Z_l(F_{X_{\epsilon, \hat{X}_{\epsilon}^*}}^{\phi_s})}_{\text{a constant}} \stackrel{\text{dist.}}{=} \underbrace{\frac{1}{2} \log \left(\frac{\det(C_{\mathbf{X}_{\epsilon, \phi_s}^l})}{\det(C_{\mathbf{S}_{\epsilon, \phi_s}^{l,*}})} \right)}_{\text{a constant}} + \underbrace{\frac{1}{2} \log(e) \cdot \tilde{V}_{\epsilon, \phi_s}^l}_{\text{a RV}}.$$

- 2 Rewrite $\tilde{V}_{\epsilon, \phi_s}^l \stackrel{\text{dist.}}{=} \sum_{i=0}^{2l - \tilde{l}_{\epsilon, l} - 1} \lambda_{\epsilon, i}^{\phi_s} \cdot \underline{\underline{(\tilde{\gamma}_{\epsilon, i}^{\phi_s})^2}}$, where $2l - \tilde{l}_{\epsilon, l} - 1 \in \mathcal{N}$, $\lambda_{\epsilon, i}^{\phi_s} \in \mathcal{R}$ and $\underline{\underline{(\tilde{\gamma}_{\epsilon, i}^{\phi_s})^2}}$ are **mutually independent central chi-square RVs** with **single degrees of freedom**.

Main Theorem 1

Proof Sketch

- ③ Calculate $\mathbb{E}\{\underline{\tilde{V}}_{\epsilon, \phi_s}^I\}$ and $\text{Var}\{\underline{\tilde{V}}_{\epsilon, \phi_s}^I\}$, then $\mathbb{E}\{\underline{Z}_I(F_{X_\epsilon, \hat{X}_\epsilon^*}^{\phi_s})\}$ and $\text{Var}\{\underline{Z}_I(F_{X_\epsilon, \hat{X}_\epsilon^*}^{\phi_s})\}$ are obtained.
- ④ By the **Chebyshev inequality** and the **limit-superior in probability**, we obtain

$$R_\epsilon^{\phi_s}(D) \leq \limsup_{l \rightarrow \infty} \min_{\mathcal{C}_{S^l}^{\epsilon, \phi_s}} \min_{\mathcal{C}_{S^l}^{\epsilon, \phi_s}} \frac{1}{2l} \log \left(\frac{\det(\mathbf{C}_{\mathbf{X}^l_{\epsilon, \phi_s}})}{\det(\mathbf{C}_{\mathbf{S}^l_{\epsilon, \phi_s}})} \right).$$

- Combining the **converse** and the **achievability**, we obtain

$$R_{\epsilon}^{\phi_s}(D) = \limsup_{l \rightarrow \infty} \min_{\mathcal{C}_{S_{\epsilon, \phi_s}^l} \in \mathcal{C}_{S_{\epsilon, \phi_s}^l}} \frac{1}{2l} \log \left(\frac{\det(\mathbf{C}_{\mathbf{X}_{\epsilon, \phi_s}^l})}{\det(\mathbf{C}_{S_{\epsilon, \phi_s}^l})} \right).$$

- Considering the **continuousness** of the source sequence generation and the **equidistribution theorem** [Kuipers and Niederreiter: 1974], [Coppel: 2009],
 - as the number of source sequences goes to **infinity**, the **sampling phase** of each source sequence becomes **asymptotically uniformly distributed** over $[0, T_c)$.
 - Therefore, the **RDF** for compressing $\mathbf{X}_{\epsilon}^{\phi_s}[i]$ is

$$R_{\epsilon}(D) = \frac{1}{T_c} \int_{\phi_s=0}^{T_c} R_{\epsilon}^{\phi_s}(D) d\phi_s.$$

• Remark 1:

- Intuitively, as $\epsilon \notin \mathcal{Q}$, the AF of $X_\epsilon^{\phi_s}[i]$ **never repeats itself** and **asymptotically exhibits all values** of $c_{X_c}(t, \lambda)$ as the **blocklength** l goes to **infinity**.
- Therefore, we conjecture that $R_\epsilon^{\phi_s}(D)$ should **NOT** be affected by the **initial sampling phase** $\phi_s \in [0, T_c)$, therefore

$$R_\epsilon^{\phi_s}(D) = R_\epsilon^0(D).$$

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$$R_\epsilon^{\phi_s}(D) = R_\epsilon^0(D).$$

• Remark 2:

- When the sampling is **synchronous** ($\epsilon \in \mathcal{Q}$), the **RDF** is **affected** by the **initial sampling phase** $\phi_s \in [0, T_c)$, namely

$$R_\epsilon^{\text{sync}}(D) = R_\epsilon^{\phi_s}(D),$$

which goes back to the result in [Kipnis et al.: *IEEE TIT* 2018].

- In this case, if the **RDF** has to be **fixed** and **minimized**, we have

$$R_\epsilon^{\text{sync}}(D) = \min_{\phi_s \in [0, T_c)} R_\epsilon^{\phi_s}(D).$$

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 - Cyclostationary Processes
 - Statistics of Sampled CT WSCS Processes
 - Rate-Distortion Theory
 - Information-Spectrum Framework
- 2 Motivating Example and Related Works
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- 3 Problem Formulation
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- 6 Summary and Future Works

Main Theorem 2

Auxiliary Theorem: The RDF for DT WSCS Processes

Auxiliary Theorem: The **RDF** for **DT WSCS Processes**:

- ① Let the **source process** $X_c(t)$ be sampled with **interval** $T_s(\epsilon_n) \triangleq \frac{T_c}{p+\epsilon_n}$, $\epsilon_n \triangleq \frac{\lfloor n \cdot \epsilon \rfloor}{n} \in \mathcal{Q}$, $n \in \mathcal{N}^+$, and with an **initial sampling phase** $\phi_s \in [0, T_c)$:

$$X_{\epsilon_n}^{\phi_s}[i] \triangleq X_c(i \cdot T_s(\epsilon_n) + \phi_s).$$

$X_{\epsilon_n}^{\phi_s}[i]$ is a **DT WSCS process** with

- maximal autocorrelation length $\tau_c \triangleq \left\lceil \frac{(p+1) \cdot \lambda_c}{T_c} \right\rceil \geq \left\lceil \frac{(p+\epsilon_n) \cdot \lambda_c}{T_c} \right\rceil$;
- period of statistics $p_n \triangleq p \cdot n + \lfloor n \cdot \epsilon \rfloor$.

Main Theorem 2

Auxiliary Theorem: The RDF for DT WSCS Processes

- Transform $X_{\epsilon_n}^{\phi_s}[i]$ into a **p_n -dimensional stationary process**:

$$\left(\mathbf{X}_{\epsilon_n, \phi_s}^{p_n}[i] \right)_m \triangleq X_{\epsilon_n}^{\phi_s}[i \cdot p_n + m], \quad m = 0, 1, \dots, p_n - 1.$$

- Autocorrelation matrix**

$$\mathbf{C}_{\mathbf{X}_{\epsilon_n, \phi_s}^{p_n}}[\Delta] \triangleq \mathbb{E} \left\{ \mathbf{X}_{\epsilon_n, \phi_s}^{p_n}[i] \cdot \left(\mathbf{X}_{\epsilon_n, \phi_s}^{p_n}[i + \Delta] \right)^T \right\}, \quad \Delta \in \mathcal{Z};$$

- Power spectral density (PSD) matrix**

$$\mathbf{S}_{\mathbf{X}_{\epsilon_n, \phi_s}^{p_n}}(f) \triangleq \sum_{\Delta \in \mathcal{Z}} \mathbf{C}_{\mathbf{X}_{\epsilon_n, \phi_s}^{p_n}}[\Delta] \cdot e^{-j2\pi f \Delta}, \quad -\frac{1}{2} \leq f \leq \frac{1}{2}.$$

- Let $\underline{\lambda}_{\epsilon_n, \phi_s, m}^{p_n}(f)$, $0 \leq m \leq p_n - 1$, denote the **eigenvalues** of $\mathbf{S}_{\mathbf{X}_{\epsilon_n, \phi_s}^{p_n}}(f)$.

Main Theorem 2

Auxiliary Theorem: The RDF for DT WSCS Processes

- ③ Given the **distortion constraint** D , the **RDF** for compressing $X_{\epsilon_n}^{\phi_s}[i]$ is [Kipnis et al.: *IEEE TIT* 2018]

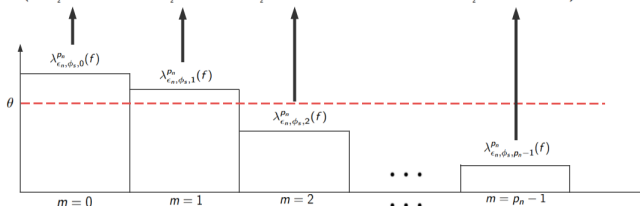
$$R_{\epsilon_n}^{\phi_s}(D) = \frac{1}{2p_n} \sum_{m=0}^{p_n-1} \int_{f=-\frac{1}{2}}^{\frac{1}{2}} \left(\log \left(\frac{\lambda_{\epsilon_n, \phi_s, m}^{p_n}(f)}{\theta} \right) \right)^+ df,$$

where θ is selected s.t.

$$D = \frac{1}{p_n} \sum_{m=0}^{p_n-1} \int_{f=-\frac{1}{2}}^{\frac{1}{2}} \min \left\{ \lambda_{\epsilon_n, \phi_s, m}^{p_n}(f), \theta \right\} df.$$

Reverse waterfilling

$$\frac{1}{p_n} \left(\int_{f=-\frac{1}{2}}^{\frac{1}{2}} \theta df + \int_{f=-\frac{1}{2}}^{\frac{1}{2}} \theta df + \int_{f=-\frac{1}{2}}^{\frac{1}{2}} \lambda_{\epsilon_n, \phi_s, 2}^{p_n}(f) df + \dots + \int_{f=-\frac{1}{2}}^{\frac{1}{2}} \lambda_{\epsilon_n, \phi_s, p_n-1}^{p_n}(f) df \right) = D$$



Main Theorem 2

The RDF with a Finite and Bounded Delay between Sequences

RDF characterization for $X_\epsilon^{\phi_s}[j]$:

1 Define $R_{\epsilon_n}(D) \triangleq \min_{\phi_s \in [0, T_c)} R_{\epsilon_n}^{\phi_s}(D)$.

2 If the **AF** of $X_c(t)$ satisfies

$$\min_{0 \leq t < T_c} \left\{ c_{X_c}(t, 0) - 2 \cdot \tau_c \cdot \max_{|\lambda| > \frac{T_c}{p+1}} \left\{ |c_{X_c}(t, \lambda)| \right\} \right\} \geq \gamma_c > 0,$$

3 for the **distortion constraint** $D \leq \gamma_c$,

4 when a **finite and bounded delay** between processed sequences up to $\tau_c \cdot T_s(\epsilon) + T_c$ in CT is allowed,

the **RDF** for compressing $X_\epsilon^{\phi_s}[j]$ is given as

$$R_\epsilon(D) = \limsup_{n \rightarrow \infty} R_{\epsilon_n}(D).$$

- **Gaussian asymptotically wide-sense stationary (AWSS) vector processes** [Gutiérrez-Gutiérrez et al.: *Entropy* 2018]
 - Such processes have a **constant mean** and their $n \times n$ **autocorrelation matrices** are **asymptotically block Toeplitz**.
 - **DT WSACS processes CANNOT** be transformed into **AWSS vector processes**.

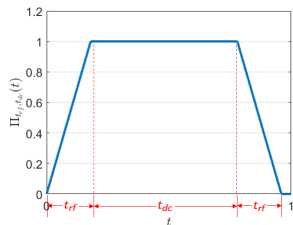
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 - **DT WSACS processes** **CANNOT** be transformed into **AWSS vector processes**.
- **AMS processes** [Gray: 2011]
 - **DT WSACS processes** are **AMS** [Gardner: *J. Sound Vib.* 2023].
 - For an **AMS process** with **reference letter**, its **DRF** is **equivalent** to the **DRF** of its **associated stationary source**.
 - In our scenario, we were **NOT** able to evaluate the AMS **stationary mean**, i.e., **probability measure** of a **Gaussian mixture** as the number of summands goes to **infinity**.

Main Theorem 2

Numerical Evaluations and Implications

Setup of CT AF:

- 1 Consider a **periodic pulse function** $\Pi_{t_{rf}, t_{dc}}(t)$, $t_{rf} = 0.01$, $t_{dc} \in [0, 0.98]$ and **period** 1:
- 2 Set **period of AF** of $X_c(t)$ as $T_c = 5 \mu\text{sec}$.

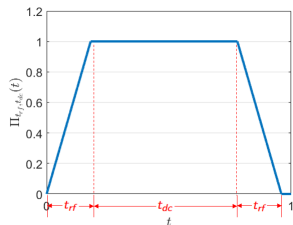


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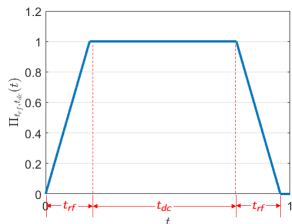
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Numerical Evaluations and Implications

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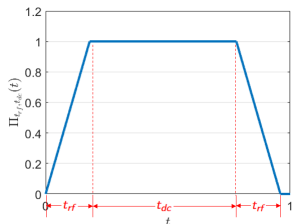
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- 4 The variance function of $X_c(t)$ is given as

$$c_{X_c}(t, 0) \triangleq 2 + 8 \cdot \Pi_{t_{rf}, t_{dc}} \left(\frac{t}{T_c} - \phi \right).$$



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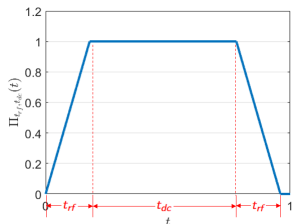
$$c_{X_c}(t, 0) \triangleq 2 + 8 \cdot \Pi_{t_{rf}, t_{dc}} \left(\frac{t}{T_c} - \phi \right).$$

- 5 Set **maximal autocorrelation length** as $\lambda_c = 4 \mu\text{sec}$. Denote the **indicator function** as $\mathbf{1}(\cdot)$. For $\lambda \geq 0$, the AF is given as

$$c_{X_c}(t, \lambda) \triangleq \mathbf{1} \left(\lambda \in [0, \lambda_c] \right) \cdot e^{-\lambda \cdot 10^{6.1}} \cdot c_{X_c}(t, 0).$$

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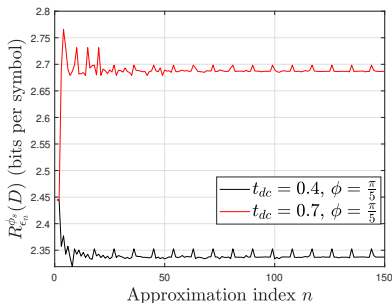
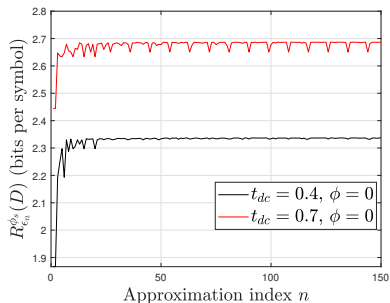
$$c_{X_c}(t, \lambda) \triangleq \mathbf{1} \left(\lambda \in [0, \lambda_c] \right) \cdot e^{-\lambda \cdot 10^6 \cdot 1} \cdot c_{X_c}(t, 0).$$

- 6 For $\lambda < 0$, $c_{X_c}(t, \lambda) = c_{X_c}(t + \lambda, -\lambda)$.

Main Theorem 2

Numerical Evaluations and Implications

The **RDF** $R_{\epsilon_n}(D)$ w.r.t. **approximation index** n for $\epsilon = \frac{\pi}{7}$:

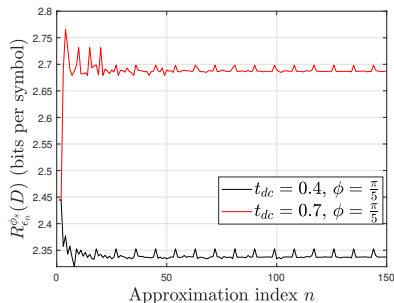
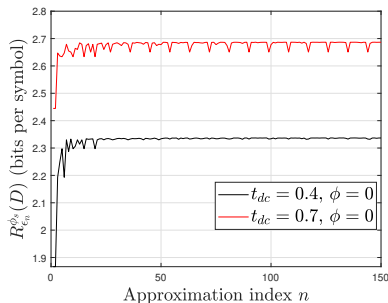


- When t_{dc} is **higher**, $R_{\epsilon_n}(D)$ is **higher**.
 - Larger **variance** \longrightarrow larger **dynamic range** \longrightarrow more **bits per symbol**.

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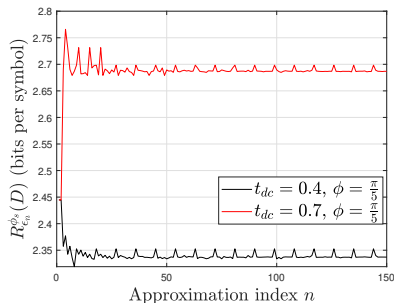
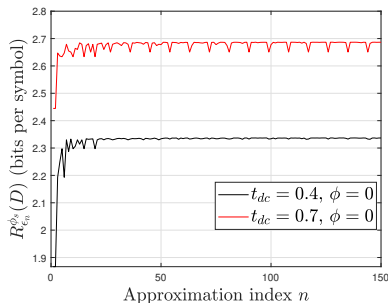


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 - Larger **variance** \rightarrow larger **dynamic range** \rightarrow more **bits per symbol**.
- When n is **large enough**, $R_{\epsilon_n}(D)$ becomes **periodically stable**.
 - Larger n \rightarrow larger **period of statistics** p_n \rightarrow shifting of sampling instances **negligible**.

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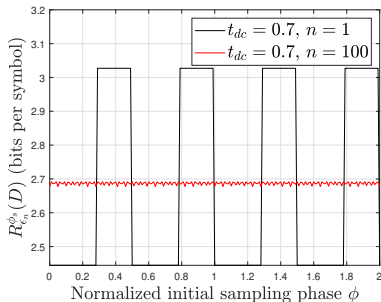
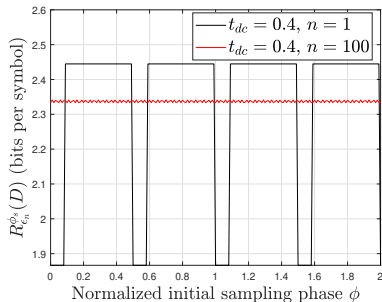


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- When n is **large enough**, $R_{\epsilon_n}(D)$ becomes **periodically stable**.
 - Larger n \rightarrow larger **period of statistics** p_n \rightarrow shifting of sampling instances **negligible**.
- $R_{\epsilon_n}(D)$ does **NOT** converge.
 - **Asynchronism** of sampling \rightarrow **nonstationarity**.

Main Theorem 2

Numerical Evaluations and Implications

The **RDF** w.r.t. **normalized initial sampling phase** ϕ for $\epsilon = \frac{\pi}{7}$:

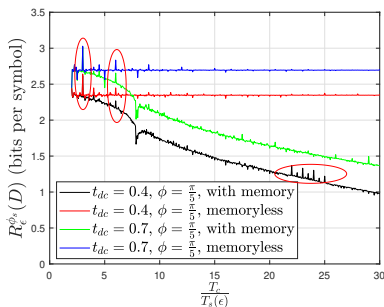
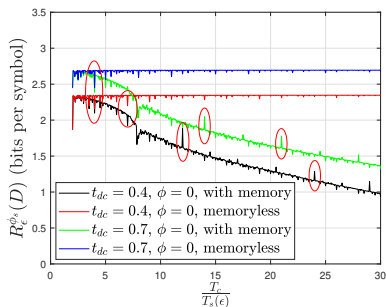


- When n is **large enough**, $R_{\epsilon_n}(D)$ varies **very little** w.r.t. **normalized initial sampling phase**
 - Larger n \longrightarrow more **CT cycles** covered in **one DT period** \longrightarrow impact of **shifting** of a few samples on the rate **negligible**.

Main Theorem 2

Numerical Evaluations and Implications

The **RDF** w.r.t. **sampling frequency**:

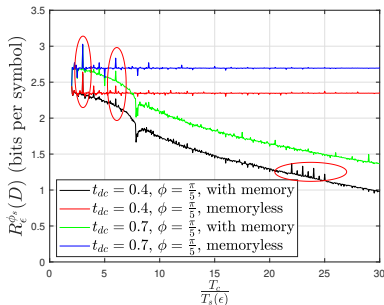
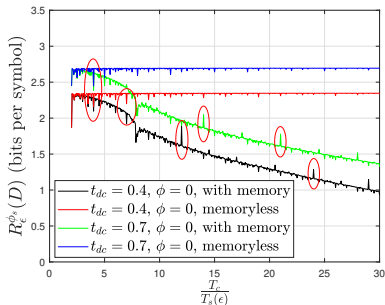


- RDF for **processes with memory** is generally **smaller** than RDF for **memoryless processes**.

Main Theorem 2

Numerical Evaluations and Implications

The **RDF** w.r.t. **sampling frequency**:



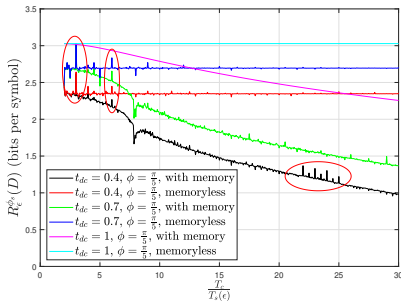
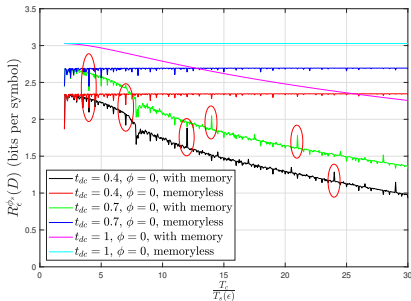
- RDF for **processes with memory** is generally **smaller** than RDF for **memoryless processes**.
- **Small sampling rate variations** could lead to **large compression rate variations**.

Practical insight: In design of communications systems, assume **asynchronous sampling**.

Main Theorem 2

Numerical Evaluations and Implications

RDF for sampled **WSCS** processes vs **RDF** for sampled **stationary** processes

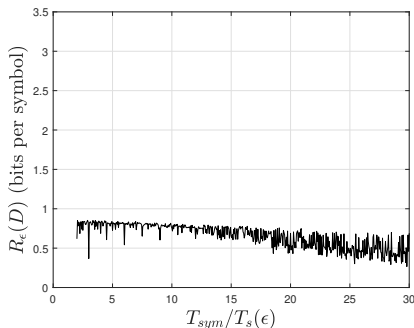


- **RDF curve** for the **stationary** sampled process is **smooth**.
- **RDF curve** for the **memoryless stationary** sampled process is **flat**.
 - Compression of an **i.i.d.** source.

Main Theorem 2

Numerical Evaluations and Implications

The **RDF** of **real passband OFDM signal** w.r.t. **sampling frequency**:



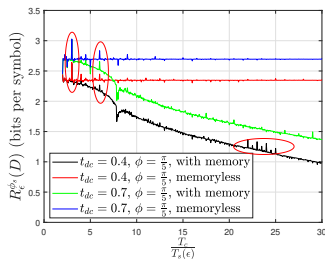
- **Real passband OFDM signal** is a **CT WSCS** process
 - with the **period of statistics** equaling to its **symbol period**.
- RDF exhibits **fluctuations**, demonstrating the **practical insight**.

Practical insight: In design of communications systems, assume **asynchronous sampling**.

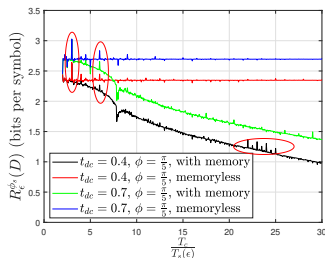
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- We defined the problem of **RDF characterization** for **DT WSACS Gaussian processes with memory**,
 - without delay between consecutive sequences.
 - with **finite and bounded processing delay** between consecutive sequences.
- This problem arises in practical sampling scenarios:
 - e.g., **compress-and-forward relay networks**.
- We presented the **RDF results** for compressing **DT WSACS Gaussian processes with memory** for two source sequence generation schemes.
- We discussed **additional** frameworks for RDF analysis and notions implied from **numerical evaluations**.
 - **Gaussian AWSS vector processes** and **AMS processes** are **NOT** useful for RDF analysis in our problem.

- **Important insight:** Assume **asynchronous sampling** in **practical communications systems design**:

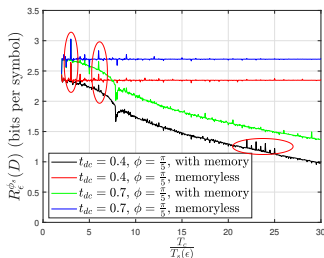


- **Important insight:** Assume **asynchronous sampling** in **practical communications systems design**:



- **Future research directions:**
 - **Compress-and-forward relaying:** Clock synchronization mismatch between the relay and the source and between the relay and the destination.

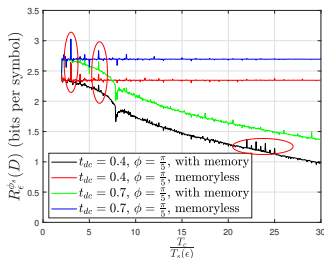
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- **Compress-and-forward relaying:** Clock synchronization mismatch between the relay and the source and between the relay and the destination.
- **Wyner-Ziv model** in multiuser compression, where the **sampling clock** of the **side information** is **NOT** synchronized with that of the **source**.
- **Secrecy capacity** in the presence of an **eavesdropper**
 - ▶ Multiple settings depending on which receivers (i.e., destination or eavesdropper) observes interference.